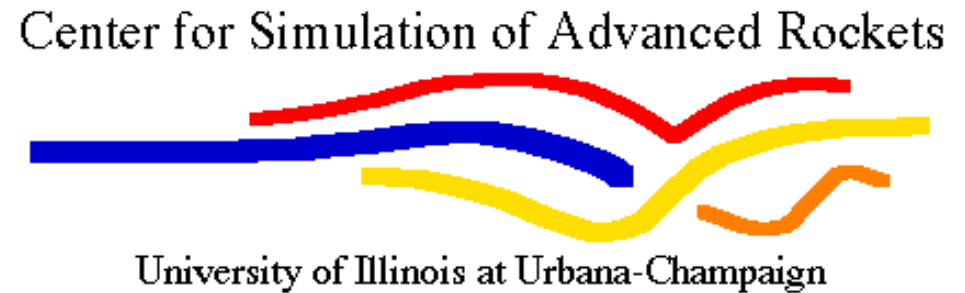


Center for Simulation of Advanced Rockets &  
Department of Computer Science  
University of Illinois at Urbana-Champaign  
<http://www.csar.uiuc.edu>



# Truncation of Optimal Krylov Subspace Methods

Eric de Sturler

✉ [sturler@uiuc.edu](mailto:sturler@uiuc.edu)



<http://www.csar.uiuc.edu/~sturler>

Sandia National Laboratories, March 15-19, 1999

## Overview

1. Truncation: why and how (not)
2. Neglecting orthogonality and generalized truncation
3. GMRES(m): Analysis and preserving a subspace
4. GCRO with optimal truncation
5. Numerical examples
6. Conclusions
7. An application in lattice QCD (multiple right hand sides)
8. An application in Jacobi-Davidson eigenvalue solvers
9. An application in matrix-free nonlinear solvers (adaptive grid pde package)

## **Truncation: why and how (not)**

- **Full GMRES/GCR too expensive in general: memory / work**
- **GMRES-like method, close to optimal with limited resources**
- **Standard strategies:      restart GMRES(m), all previous info lost  
   truncate GCR(m), oldest vectors removed**
- **VERY CRUDE!**
- **Better to make some compromise: GMRESR / GCRO / FGMRES  
eventually these methods still need some truncation**
- **How to truncate:**
  - **careful select what to keep**
  - **don't think about vectors, think about spaces**

## (truncation: why and how (not))

**How to select the subspace to keep?**

**1. Make assumptions on system (normal, nice spectrum): restrictive**

**2. No assumptions, only use information from iteration process**

- **No general convergence theory for GMRES (GCR)**
- **Worst case exists : bad spectrum  
non-normality**
- **Optimality is obtained through orthogonality: the old 'search space'  
(after making update) is used only to maintain orthogonality.**
- **So importance of a subspace is determined by its contribution to  
maintaining orthogonality**

## GCR and nested methods

**GCR:  $Ax = b$**

$$r_0 = b - Ax_0 ; i = 0$$

**while  $\|r_i\|_2 > tol$**

$$i = i+1$$

$$u_i = r_{i-1} ; c_i = Au_i$$

$$u_i = \alpha_i (u_i - U_{i-1} C_{i-1}^H c_i)$$

$$c_i = \alpha_i (I - C_{i-1} C_{i-1}^H) c_i$$

$$r_i = r_{i-1} - c_i c_i^H r_i$$

$$x_i = x_{i-1} + u_i c_i^H r_{i-1}$$

**after k iterations:**

$$range(U_k) = K^k(A, r_0)$$

$$AU_k = C_k \quad C_k^H C_k = I$$

$$r_k = (I - C_k C_k^H) r_0$$

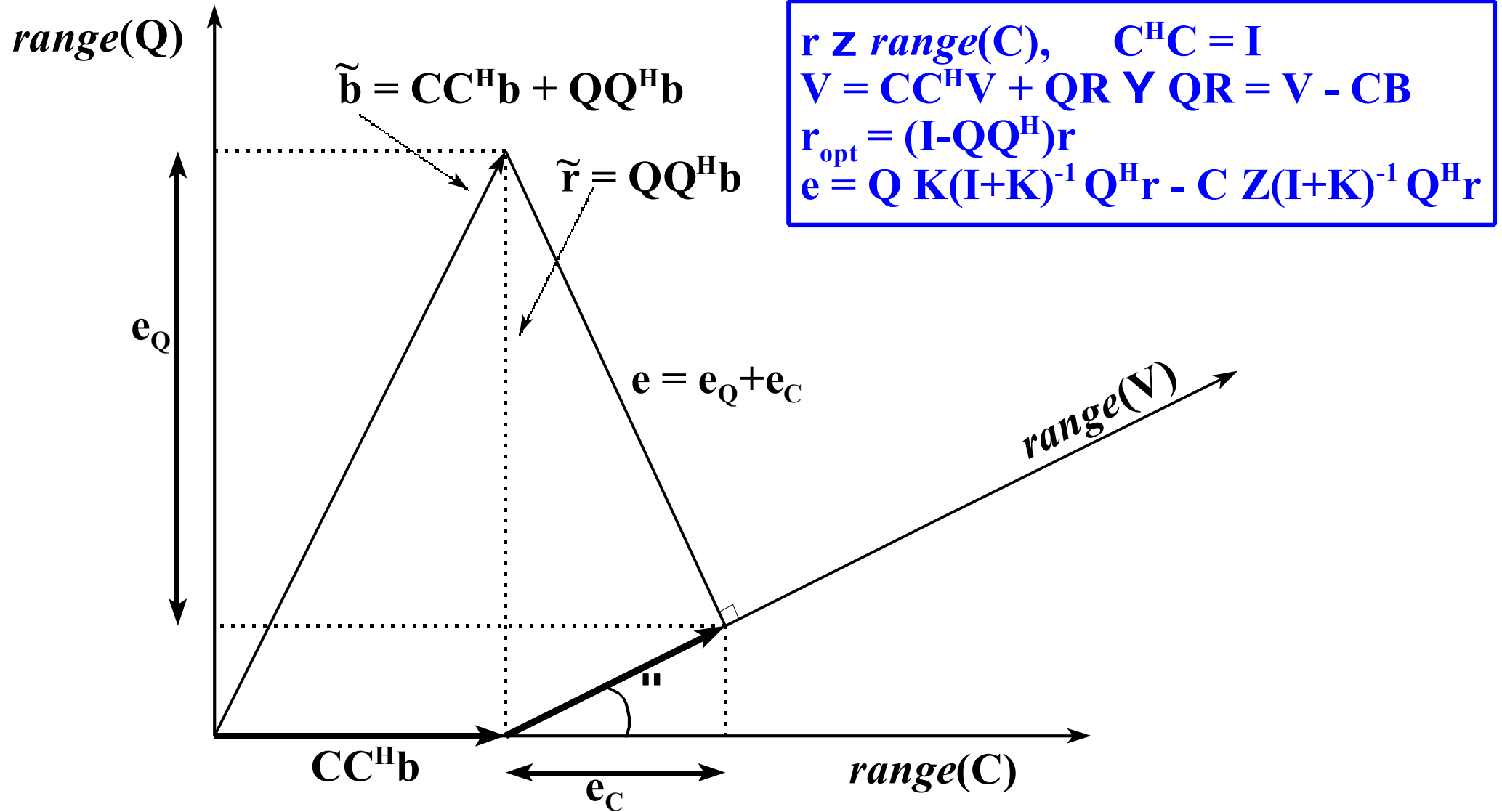
$$x_k = x_0 + U_k C_k^H r_0$$

**GMRESR: replace  $u_i = r_{i-1}$  by GMRES approx. sol. of  $Ae_{i-1} = r_{i-1}$**

**Ignores orthogonality / optimality w.r.t.  $range(C_{i-1})$**

**GCRO: maintain orthogonality; use  $(I - C_{i-1} C_{i-1}^H)A$  for inner iteration**

# Neglecting Orthogonality & Generalized Truncation



## (neglecting orthogonality & generalized truncation)

$$\mathbf{V} = \mathbf{C}\mathbf{B} + \mathbf{Q}\mathbf{R} = \mathbf{C}\mathbf{C}^H\mathbf{V} + \mathbf{Q}\mathbf{Q}^H\mathbf{V}$$

$$\mathbf{r}_{\text{opt}} = (\mathbf{I} - \mathbf{Q}\mathbf{Q}^H)\mathbf{r}$$

$$\mathbf{e} = \mathbf{Q}\mathbf{K}(\mathbf{I} + \mathbf{K})^{-1}\mathbf{Q}^H\mathbf{r} - \mathbf{C}\mathbf{Z}(\mathbf{I} + \mathbf{K})^{-1}\mathbf{Q}^H\mathbf{r}$$

$$\mathbf{Z} = \mathbf{B}\mathbf{R}^{-1} = \mathbf{X}_Z\mathbf{E}_Z\mathbf{Y}_Z^H, \quad \mathbf{K} = \mathbf{Z}^H\mathbf{Z}, \quad \langle_i = \mathbf{y}_i^H\mathbf{Q}^H\mathbf{r}$$

$$\mathbf{e} = \mathbf{E} \left( \frac{\langle_i \mathbf{F}_i^2}{1 + \mathbf{F}_i^2} \mathbf{Q}\mathbf{y}_i - \frac{\langle_i \mathbf{F}_i}{1 + \mathbf{F}_i^2} \mathbf{C}\mathbf{x}_i \right) \quad 2e2' \quad \dagger \quad \mathbf{E} \left[ \frac{\langle_i^* \mathbf{F}_i^2}{1 + \mathbf{F}_i^2} \right]^{1/2}$$

(neglecting orthogonality & generalized truncation)

$[T^*T_c]$  unitary and  $rank([T^*T_c]) = rank(C)$ ,  $\bar{V} = [CT^*V]$  and  $\bar{C} = [CT_c]$

$$\bar{Z} = \bar{B}\bar{R}^{-1} = [0^*T_c^H Z] = [0^*T_c^H X_Z G_Z Y_Z^H]$$

Select  $T_c^H$  so that singular values of  $\bar{Z}$  minimal:

$$\min_{\substack{\dim S=k-p \\ u \in S}} (\max_{u \in S} \|u^H Z\| / \|u\|) = F_{p+1} \quad \text{and} \quad S = span\{x_{p+1}, \dot{y}, x_k\}$$

Set  $T_c = [x_{p+1} \ x_{p+2} \ \dot{y} \ x_k]$   $Y T = [x_1 \ \dot{y} \ x_p]$

$$e = \sum_{i=p+1}^{\min(k,m)} \left( \frac{\langle_i F_i^2}{1+F_i^2} Qy_i - \frac{\langle_i F_i}{1+F_i^2} Cx_i \right)$$

$$2e^2 \leq \sum_{i=p+1}^{\min(k,m)} \frac{\langle_i^2 F_i^2}{1+F_i^2} \check{L}^{1/2}$$

## GMRES(m) Analysis & Selecting a Subspace

$$\text{GMRES: } \mathbf{A}\mathbf{W}_m = \mathbf{W}_{m+1}\bar{\mathbf{H}}_m = \mathbf{W}_{m+1}\mathbf{Q}_m\bar{\mathbf{R}}_m = \mathbf{W}_{m+1}\bar{\mathbf{Q}}_m\mathbf{R}_m$$

How would convergence have been after restarting after  $s < m$  iterations?

What was influence on later iterations of orthogonality against  $\text{range}(\mathbf{A}\mathbf{W}_s)$ ?

$$\mathbf{r}_s = \mathbf{W}_{s+1}\mathbf{D}_s \quad \text{with } \mathbf{D}_s = (\mathbf{I} - \bar{\mathbf{Q}}_s\bar{\mathbf{Q}}_s^H)\mathbf{r}_s$$

$$\mathbf{C}_s : \text{orthogonal basis for } \text{range}(\mathbf{A}\mathbf{W}_s) \quad ! \quad \mathbf{C}_s = \mathbf{W}_{m+1}\bar{\mathbf{Q}}_s$$

$$\mathbf{V}_s : \text{any basis for } \mathbf{A}\mathbf{K}^{m-s}(\mathbf{A}, \mathbf{r}_s) = \mathbf{W}_{m+1}\bar{\mathbf{H}}_m\mathbf{K}^{m-s}(\bar{\mathbf{H}}_m, \mathbf{D}_s) = \tilde{\mathbf{V}}_s$$

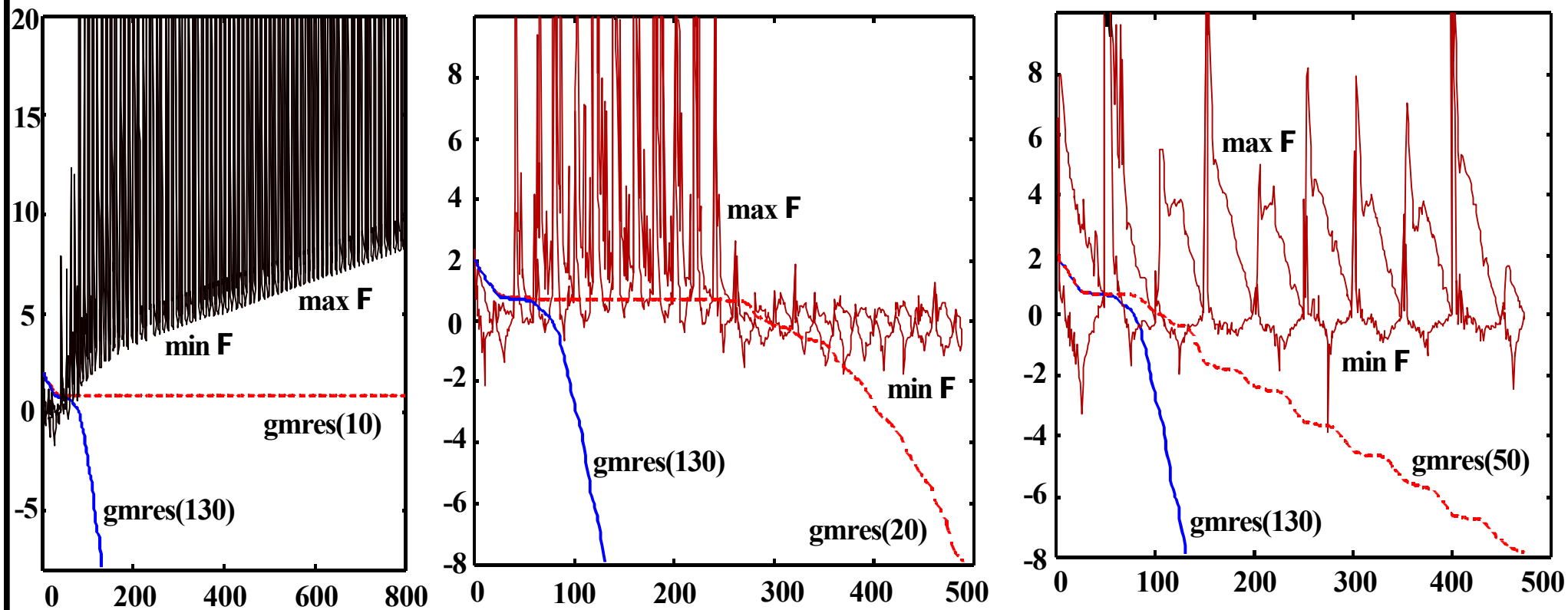
$\mathbf{Q}_s$  : orthogonal complement of  $\mathbf{C}_s$  in  $\text{range}(\mathbf{A}\mathbf{W}_m)$

$$[\mathbf{C}_s | \mathbf{Q}_s] = \mathbf{W}_{m+1}\bar{\mathbf{Q}}_m$$

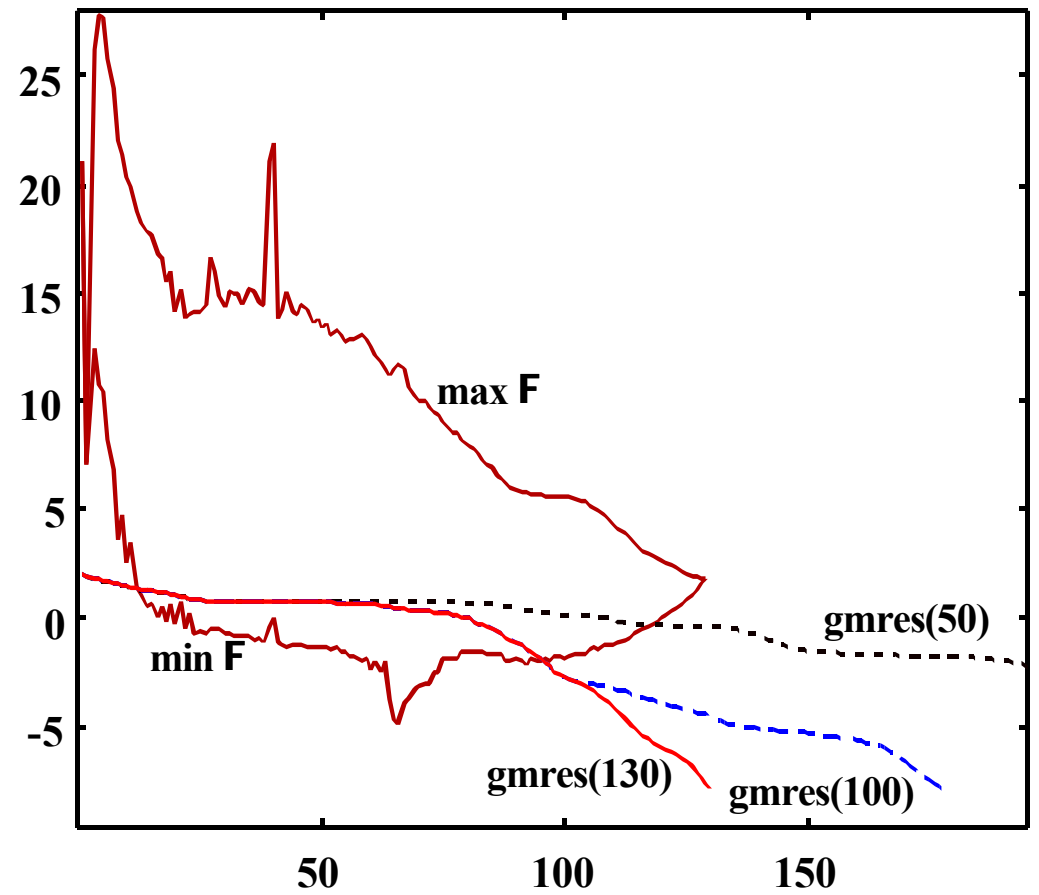
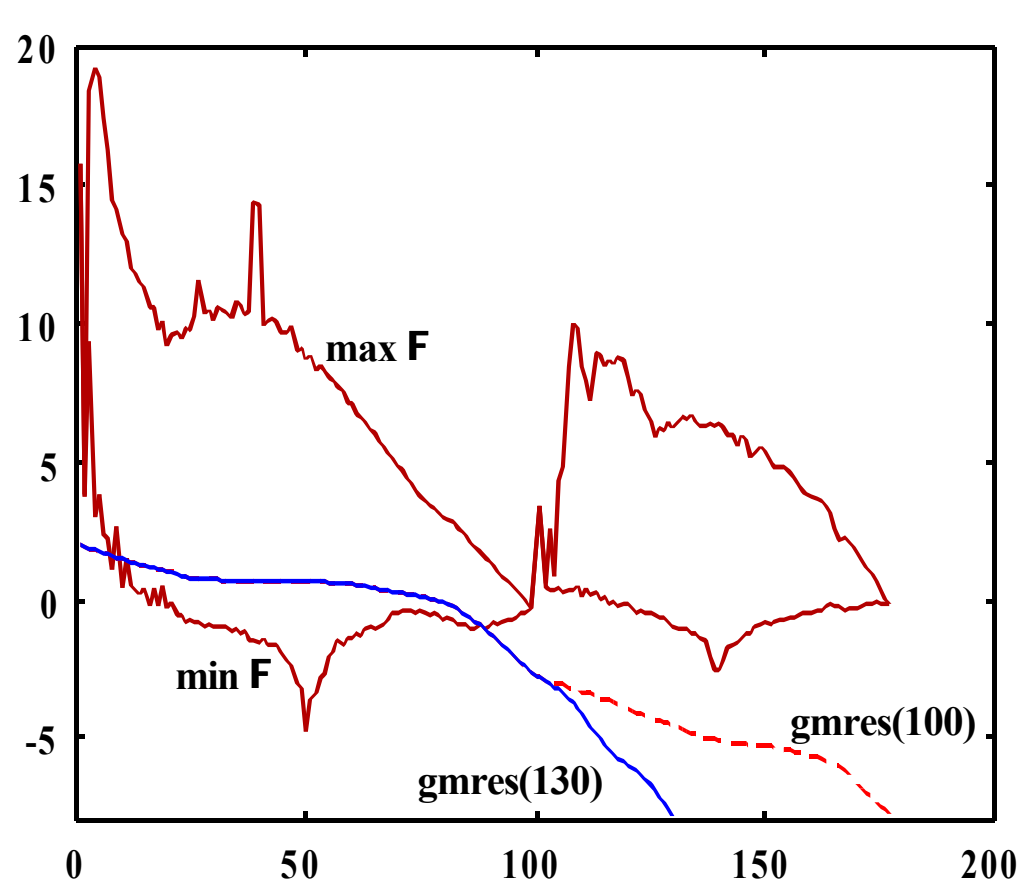
$$\bar{\mathbf{Q}}_m^H \tilde{\mathbf{V}}_s = \begin{matrix} \mathbb{R} & \mathbb{B}_s & \mathbb{X} \\ \mathbb{M} & \mathbb{R}_s & \end{matrix} \quad ! \quad \mathbf{Z}_s$$

1. difference between optimal residual  $\mathbf{r}_m$  and residual after restart
2. important vectors for setting  $\mathbf{V}$  'orthogonal' to  $\mathbf{C}$  (useful to keep)

# (GMRES(m) analysis & selecting a subspace)



# (GMRES(m) analysis & selecting a subspace)



## GCRO with Optimal Truncation

Outer method:  $U_k, C_k$  such that

$$\begin{aligned} AU_k &= C_k & C_k^H C_k &= I_k & (\text{cf. GCR}) \\ \mathbf{r}_k &= (\mathbf{I} - C_k C_k^H) \mathbf{b} & \mathbf{x}_k &= U_k C_k^H \mathbf{b} \end{aligned}$$

Inner GMRES (optimal solution over  $\text{range}(U_k)$   $\mathbf{r}$   $\text{range}(V_m)$ ):

$$\begin{aligned} AV_m &= C_k B_m + V_{m+1} H_m & \text{where } B_m &= C_k^H AV_m \\ (\mathbf{I} - C_k C_k^H) AV_m &= V_{m+1} H_m \end{aligned}$$

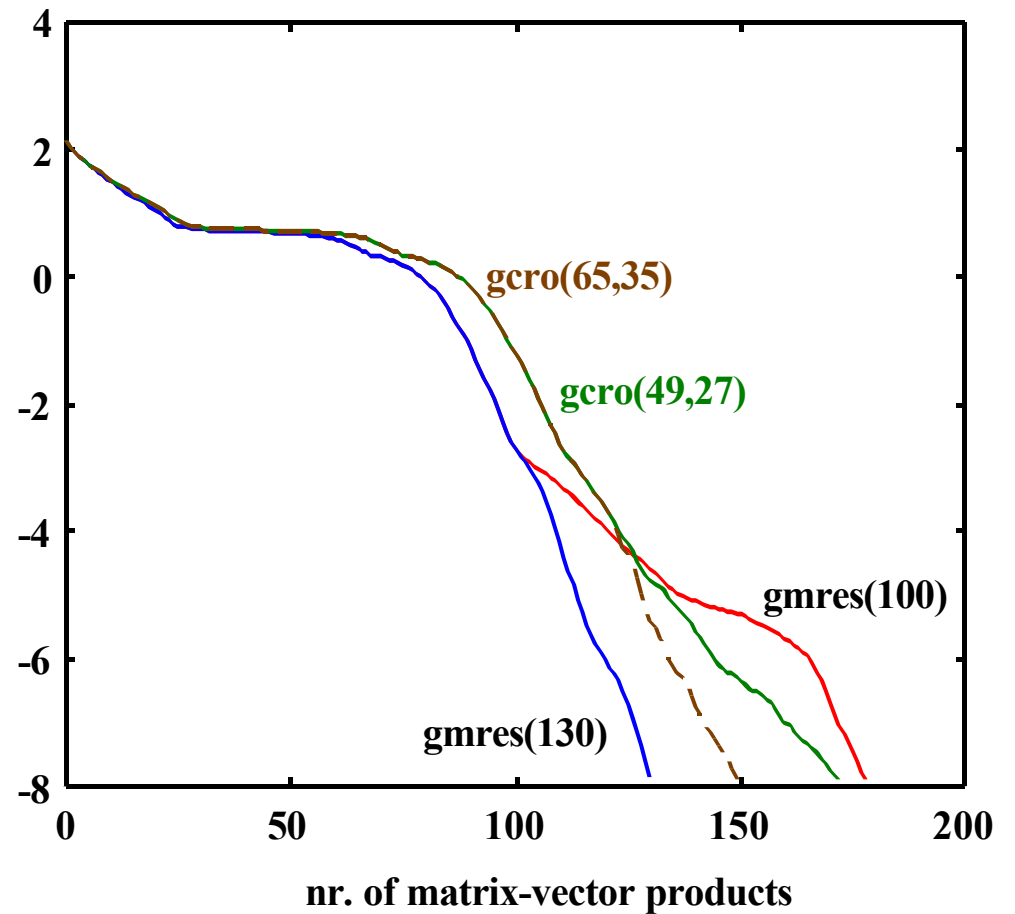
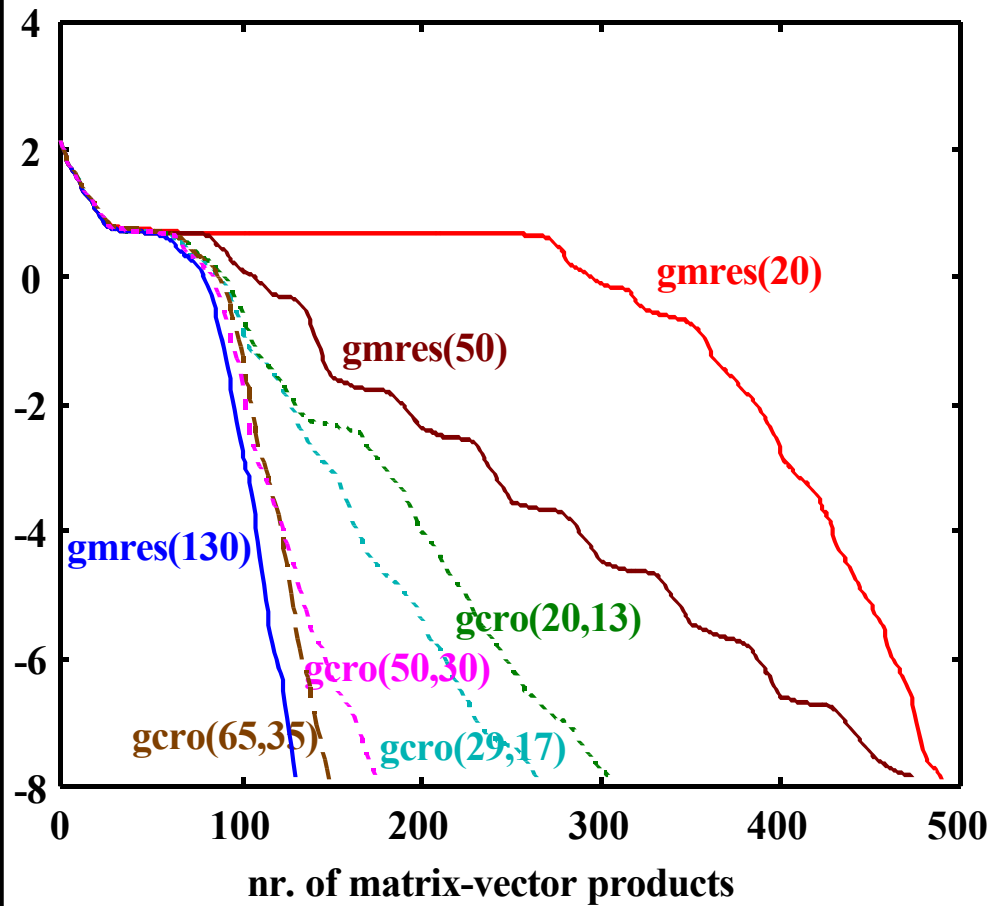
$$\mathbf{y} = \arg \min^{**} \mathbf{r}_k - V_{m+1} H_m \mathbf{y}^{**}$$

$$\mathbf{r}_{k+1} = \mathbf{r}_k - V_{m+1} H_m \mathbf{y}$$

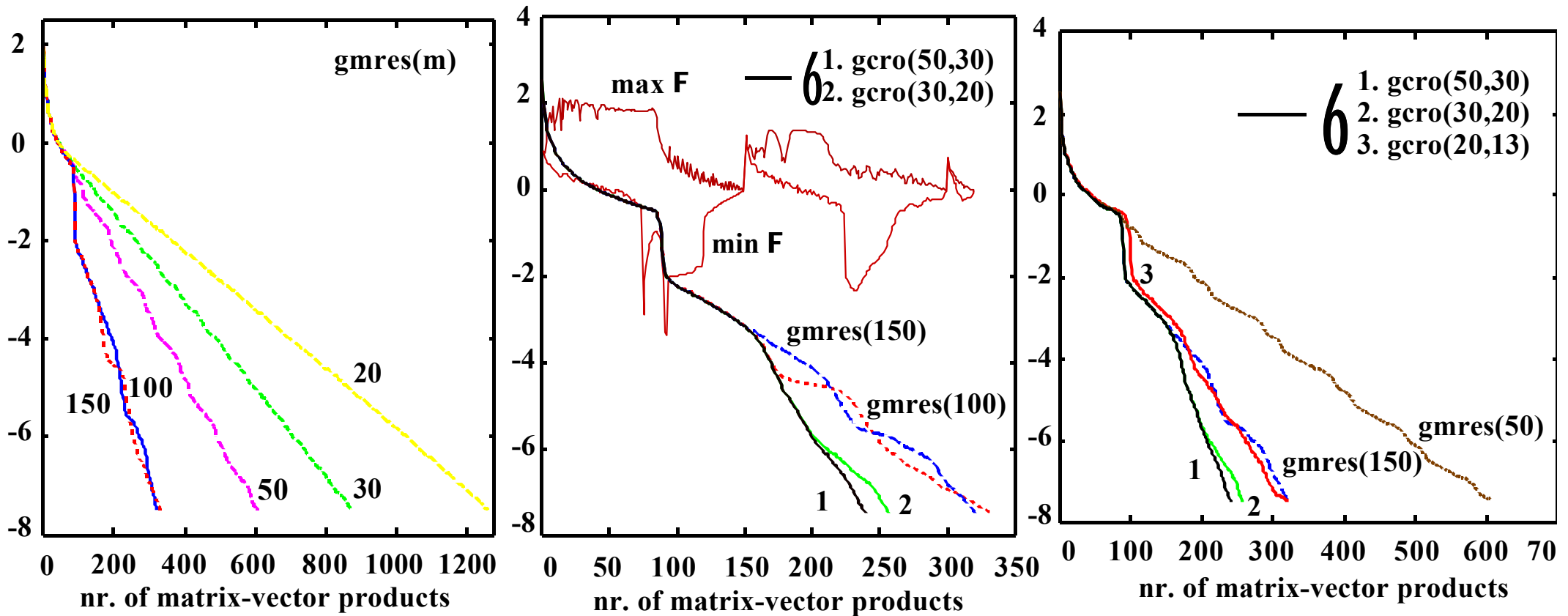
$$\mathbf{c}_{k+1} = V_{m+1} H_m \mathbf{y} = (\mathbf{I} - C_k C_k^H) AV_m \mathbf{y}$$

$$\mathbf{u}_{k+1} = A^{-1} V_{m+1} H_m \mathbf{y} = A^{-1} (\mathbf{I} - C_k C_k^H) AV_m \mathbf{y} = V_m \mathbf{y} - U_k C_k^H AV_m \mathbf{y}$$

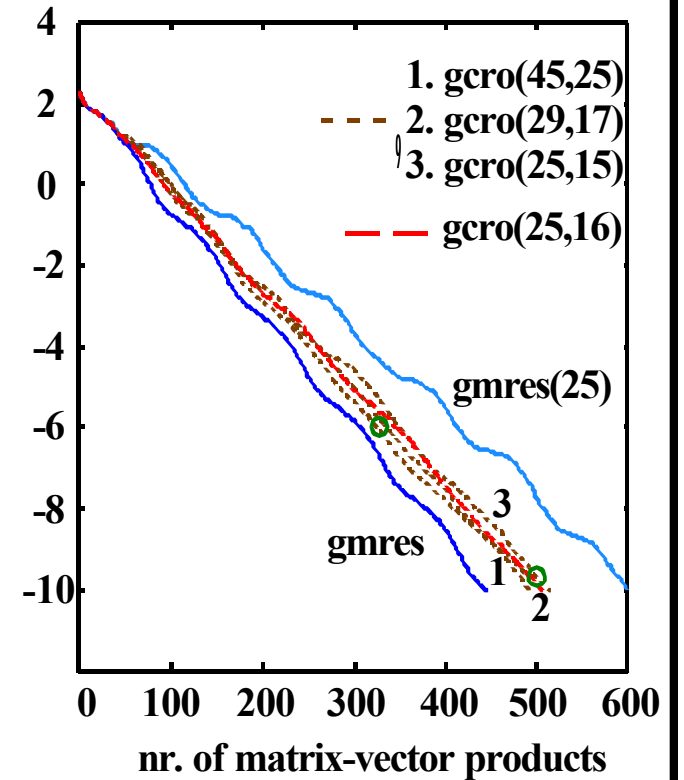
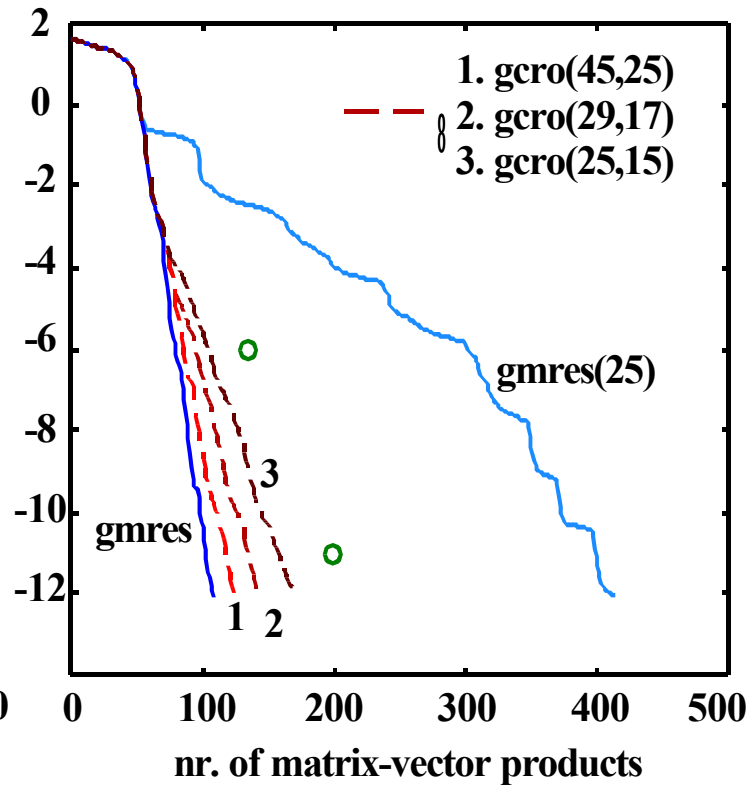
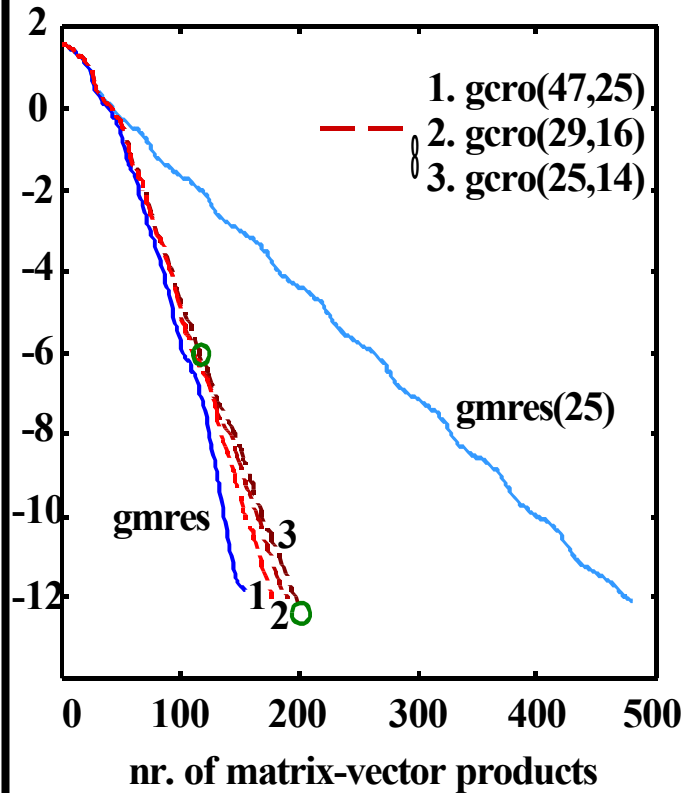
# (GCRO with optimal truncation)



# (GCRO with optimal truncation)



# (GCRO with optimal truncation)

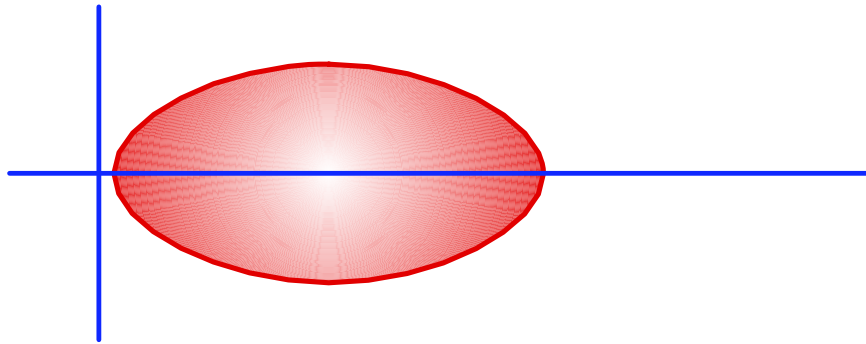


## **Conclusions and Future Work**

- **Almost optimal convergence with very small number of vectors**
- **Minimizing over selected subspaces of very small dimension works better than GMRES(m) with large m**
- **Methods seems to select right subspaces to keep**
- **Initial tests indicate comparable or better than Morgan's Augm. GMRES**
- **Method insensitive to small changes in parameters, and gradual improvement for increase in number of vectors**
- **Theoretical implications / Analysis**
- **More flexible and adaptive algorithm; theory / heuristics**

# Lattice QCD

- 4D Space-time lattice, particles move in 4 directions and change color (3 colors)
- Matrix looks like a 4D Laplacian where each off-diagonal element is replaced by 12x12 (random) unitary matrix (complex).
- These unitary matrices are different for each off-diagonal element: no smoothness.
- $A = (I - kM)$ , with parameter  $k$ ; most interesting case when  $A$  almost singular.
- Eigenvalues of  $M$  uniformly distributed in (known) ellipse (depending on other parm).
- We need to solve for multiples of 12 right hand sides (sometimes consecutive).



Use  $U$  and  $C$  matrices (approx. inverse) from previous iteration(s) to speed up convergence.

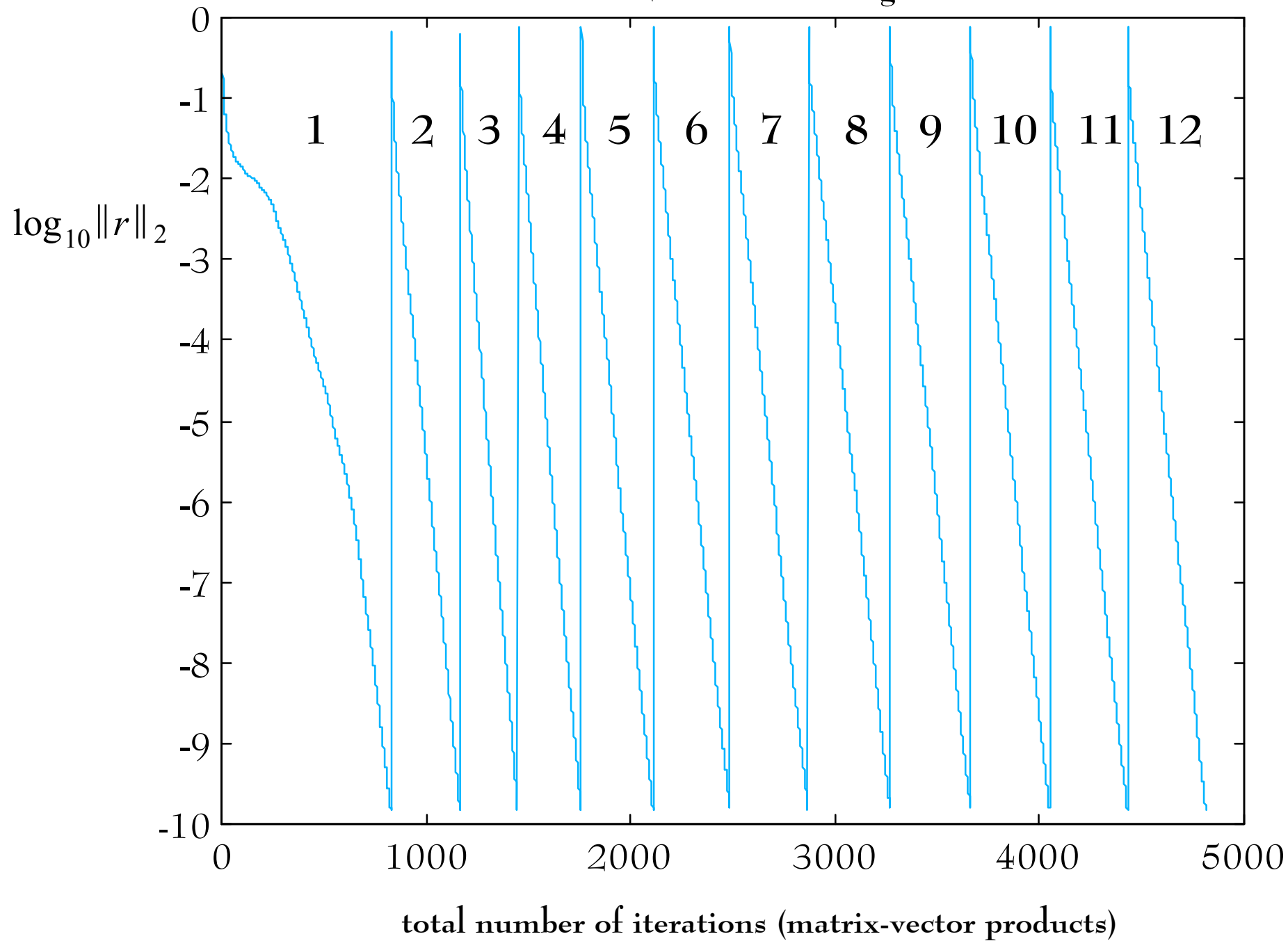
To solve  $Ax^{(k)} = b^{(k)}$  we set

$$x_0 = UC^*b \text{ and } r_0 = (I - CC^*)b$$

Iterate with method described (continuing truncation).

# Lattice QCD

GCROT for 12 consecutive right hand sides



## Jacobi-Davidson

Given  $V_k$ ;  $V_k^*V_k = I_k$ ;  $W_k = AV_k$ ;  $H_k = V_k^*W_k$  **6** eigenpair  $(q_k, s_k)$  **6**

Ritz pair  $(q_k, u_k)$ ,  $u_k = V_k s_k$ :  $r_k = Au_k - q_k u_k = (A - q_k I)u_k$

Extend the search space and compute better approximate eigenpair

Solve:  $(I - u_k u_k^*)(A - q_k I)(I - u_k u_k^*)t = -r$ ,  $t \perp u_k$

$t = (I - V_k V_k^*)t / \sqrt{2(I - V_k V_k^*)t^2}$ ;  $V_{k+1} = [V_k | t]$ ;  $W_{k+1} = [W_k | At]$ ;  $H_{k+1} = V_{k+1}^* W_{k+1}$

- Condition number of  $(I - u_k u_k^*)(A - q_k I)(I - u_k u_k^*)$ ?
- Angle  $\mathbf{p}(t, \text{range}(V_k S_c))$ ? Extension of space  $\text{range}(V_k)$ ?
- Repeatedly solving 'almost' same system (change  $(q_k, u_k)$ )
- Efficiency/Stability

# Jacobi-Davidson

